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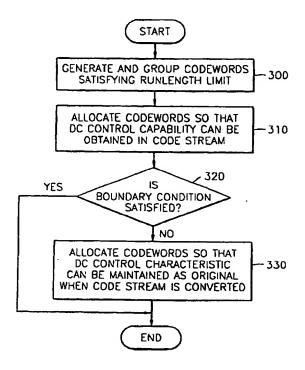
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(54) Code generation and allocation method

(57)A method for generating and allocating codewords is provided. The method includes allocating one of two selectable codewords b1 and b2 as codeword b when a preceding codeword a and a following codeword b form a code stream X, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords. According to the method, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density. Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

FIG. 3



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Description

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[0001] The present invention relates to generation and allocation of modulation codes of source codes to be recorded on a recording medium, and more particularly, to a method for generating and allocating codewords in which codewords having a restricted run length are generated and the generated codewords are allocated so that the DC control characteristic of a code stream is maintained.

[0002] In a Run Length Limited (RLL) code represented by (d, k, m, n), the performance of a code is evaluated mainly based on the recording density and the capability to suppress the DC component of the code. Here, m denotes the number of data bits (the number of so-called source data bits, which is also referred to as the number of information word bits), n denotes the number of codeword bits after modulation (the number of so-called channel bits), d denotes the minimum number of a series of '0s' that can exist between '1' and '1' in a codeword, and k denotes the maximum number of a series of '0s' that can exist between '1' and '1' in a codeword. The interval between bits in a codeword is represented by T.

[0003] In a modulation method, to improve recording density it is used to reduce the number of codeword bits n while regarding d and m as given conditions. In the RLL code, however, d which is the minimum number of a series of '0s' that can exist between '1' and '1' in a codeword, and k which is the maximum number of a series of '0s' that can exist between '1' and '1' in a codeword, should be satisfied. If, with this (d, k) condition satisfied, the number of data bits is m, the number of codewords satisfying RLL(d, k) should be equal to or greater than 2^m. Moreover, in order to actually use this code, run length constraints, that is, RLL(d, k) conditions, should be satisfied in a part where a codeword is linked to another codeword. In addition, when the DC component of a code affects the system performance, it is desirable to use a code which has a DC suppression capability.

[0004] The main reason for suppressing the DC component in the RLL modulated code stream is to minimize a reproducing signal's affect on a servo band. Hereinafter, methods for suppressing the DC component will be referred to as Digital Sum Value (DSV) control methods.

[0005] DSV control methods can be broadly classified into two types. One is a method having a DSV control code itself, where the DSV control code is capable of controlling a DSV. The other is a method of inserting a merge bit at each DSV control time. An Eight to Fourteen Modulation plus (EFM+) code performs DSV control using a separate code table, while an EFM code or a (1, 7) code performs DSV control by inserting a merge bit.

[0006] Therefore, the shape of the prior art modulation code group having the DSV control code itself capable of controlling suppression of the DC component and satisfying the conditions described above is as shown in Figure 1, in which each of a predetermined number of main conversion code groups has a corresponding code group for controlling suppression of the DC component. Each main conversion code group and its corresponding code group form a pair so that the DC component can be suppressed and controlled. In this case, there are some characteristics that distinguish codewords of the predetermined main conversion code groups. That is, there are no identical codewords between the main conversion code groups A and B. If duplicated codes are used, there might be the conversion code groups C and D for demodulating the duplicated codes, where there are no identical codewords between the conversion code groups C and D, but codewords in the code group A or B may be in the code group C or D for demodulating duplicated codes. The number of codewords in the main conversion code groups A and B and the conversion code groups C and D for demodulating duplicated codes is 2^m if the number of bits in the source word before conversion is m. [0007] If code groups E through H are DC suppression control code groups used for suppressing DC components together with code groups A through D, respectively, the characteristics of codewords in each of the code groups E through H are the same as the characteristics of codewords in the main code groups A through D respectively. That is, the same conditions for generating duplicated codewords or the same conditions for determining the number of lead zeros in a codeword are applied to each of the DC suppression control code groups E through H for controlling suppression of DC components and the conversion code groups A through D.

[0008] For example, the characteristics of the EFM+ code, which is used in current Digital Versatile Discs (DVD), has a run length condition of RLL(2, 10) and a codeword length (n) of 16 bits, is as shown in Figure 2. The main conversion code groups are MCG1 ("A" in Figure 1) and MCG2 ("B" in Figure 1) and the conversion code groups for demodulating duplicated codes are DCG1 ("C" in Figure 1) and DCG2 ("D" in Figure 2). There are four DSV code groups ("E~H" in Figure 1) which make pairs with respective conversion code groups to control suppression of DC components. There are no identical codewords between the four conversion code groups and the four DSV code groups which are code groups for controlling DC components.

[0009] Also, the conditions for generating duplicated codewords in the entire code groups are the same, and the characteristics of codewords in each code group pair that can control DC components (MCG1 and the first DSV code group, MCG2 and the second DSV code group, DCG1 and the third DSV code group, or DCG2 and the fourth DSV code group) are the same.

[0010] That is, a codeword having a continuous sequence of from 2 to 5 zeros from the Least Significant Bit (LSB) of the codeword is generated using duplicated codewords. This rule is applied to each code group in the same manner.

In each of the codewords of the first DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the main conversion code group MCG1, there is a continuous sequence of between 2 and 9 '0s' from the Most Significant Bit (MSB). In each of the codewords of the second DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the main conversion code group MCG2, there is either 0 or 1 '0' continuing from the MSB. Some bits (here, b15(MSB) or b3) in the codewords of the third DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the conversion code group DCG1 for demodulating duplicated codes are '0b', while some bits (here, b15(MSB) or b3) in the codewords of the fourth DSV code group for controlling suppression of DC components, which controls suppression of DC components together with the conversion code group DCG2 for demodulating duplicated codes, some bits (here, b15(MSB) and b3) are '1b". In developing 8 to 15 modulation code which has an advantage in the recording density aspect compared to the prior art modulation method EFM+ which uses the modulation code group shown in Figure 1 or 2, the original characteristics of a code stream change when a change occurs in a codeword because of a boundary rule applied to the locations adjacent to a boundary which connects a codeword to another codeword.

[0011] It is an aim of the present invention to provide a method for generating and allocating codewords in which a codeword having a run length restriction is generated and the codeword is allocated so that the original characteristics of a code stream are maintained without change even when a codeword is replaced according to the boundary rule when a code stream is allocated.

[0012] According to a first aspect of the present invention there is provided a method for generating and allocating codewords of source words which are to be recorded on a recording medium, the method including generating codewords satisfying predetermined run length conditions and grouping codewords according to each run length condition; and allocating the codewords such that a code(word) for the source word is capable of controlling suppression of DC components.

[0013] It is preferable that when a predetermined boundary condition is not satisfied in the code stream, allocating codewords such that codewords which satisfy the boundary condition and maintain the DC control characteristics which are considered when the initial codewords are allocated replace the initial codewords.

[0014] It is preferable that the step for generating codewords includes generating codewords satisfying the length of a predetermined first codeword, and predetermined run length conditions, grouping the codewords according to each predetermined run length condition to generate a main conversion codeword table; generating DC control codewords satisfying the length of a predetermined second codeword, and predetermined run length conditions in order to control DC components in the code(word) stream, grouping the DC control codewords, and to generate a code conversion table for controlling DC components; and generating additional DC control codewords by taking codewords which satisfy the predetermined run length conditions and are not needed in the main conversion codeword table, and grouping the additional DC control codewords.

[0015] According to a second aspect of the present invention there is provided an allocation method for allocating codewords generated for source words to be recording on a recording medium, the method including when a preceding codeword a and a following codeword b form a code stream X, allocating one of two selectable codewords b1 and b2 as codeword b, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords.

[0016] It is preferable that when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword a in the Most Significant Bit (MSB) direction is 0, and the number of continuous '0s' from the MSB of the codewords b1 or b2 in the LSB direction is 1, code changes of either the codeword a or b1(b2) occur to satisfy the boundary condition.

[0017] It is preferable that when the number of continuous '0s' between the codewords a and b is 1 or 0, the codeword a or b is changed such that the number of 0s forming the boundary is greater than 2 and less than 10.

[0018] It is preferable that the codeword a of the code stream X1 and the codeword a of the code stream X2 are changed to other codewords such that the resulting codewords a of code streams X1 and X2 have the same INV value, and as a result, by the INVs of codewords b1 and b2 following the codewords a respectively, the INVs of the X1 and X2 become different.

[0019] According to a third aspect of the present invention there is provided an allocation method for allocating codewords of source words to be recording on a recording medium, the method including when a preceding codeword b and a following codeword c form a code stream Y, allocating one of two selectable codewords b1 and b2 as the codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of b1 and c is Y1, and

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the code stream of b2 and c is Y2, allocating codewords such that INVs of Y1 and Y2 are maintained to be opposite when the codeword b1, b2 or c should be replaced by another codeword in compliance with a predetermined boundary condition between codewords.

[0020] It is preferable that when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s from the Least Significant Bit (LSB) of the codeword c toward the Most Significant Bit (MSB) is 1, codeword b which does not satisfy the boundary condition and is xxxxxxxxxxxx1001 or xxxxxxxxxxx10001 appears in both b1 an b2.

[0021] For a better understanding of the invention, and to show how embodiments of the same may be carried into effect, reference will now be made, by way of example, to the accompanying diagrammatic drawings in which:

Figure 1 is a diagram of an example of the shape of a prior art modulation code group;

Figure 2 is a table showing the prior art code group and the characteristics of codewords included in the code group;

Figure 3 is a flowchart showing a method for generating and allocating codes according to the present invention;

Figure 4 is a table showing a variety of codeword groups of main conversion tables and the characteristics of codewords in each codeword group;

Figure 5 is a table showing a variety of codeword groups of a conversion table for DC control and the characteristics of codewords in each codeword group;

Figure 6 is a table showing a variety of codeword groups of an auxiliary conversion table for DC control and the characteristics of codewords in each codeword group;

Figure 7 is a diagram for showing what should be considered for the run length conditions when codewords a and b are connected;

Figure 8 is a table showing an example of changes in INV before code conversion and after code conversion when the run length conditions described in Figure 7 are not satisfied;

Figure 9 is a diagram showing an example of code stream branching due to selective codewords b1 and b2 for DC control;

Figure 10 is a graph showing the relationship between INV values of a code stream pair;

Figures 11a through 11e are main conversion code according to the present invention;

Figures 12a through 12j are code conversion tables for DC control according to the present invention;

Figures 13a and 13b are auxiliary code conversion tables for DC control according to the present invention; and

[0022] Figure 14 is a graph showing the difference between the frequency spectrum when codewords of the code conversion table for DC control according to the present invention are used in 25% of all of the codewords, and the frequency spectrum when prior art EFM+ modulation codewords are used.

[0023] Figure 3 is a flowchart showing a method for generating and allocating codes according to a preferred embodiment of the present invention. According to the method for generating and allocating codewords of source words to be recorded on a recording medium, codewords satisfying predetermined run length conditions are generated and the generated codewords are grouped according to each run length condition in step 300. The codewords are allocated so that the code(word) streams for source words are capable of controlling DC components in step 310. It is determined whether or not predetermined boundary conditions are satisfied in the code stream in step 320. If the conditions are not satisfied, the codewords are replaced by codewords satisfying the boundary conditions while the DC control characteristics which are considered when the original codewords are allocated can be kept.

[0024] Code tables of the codewords for source code conversion are roughly divided into three types: 1) main conversion tables, 2) conversion tables for controlling DC components, and 3) auxiliary conversion tables for controlling DC components.

[0025] Figure 4 is a table showing a variety of codeword groups of main conversion tables and the characteristics of codewords in each code group. It is assumed that d denotes the minimum run length limit of a codeword, k denotes

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the maximum run length limit of a codeword, m denotes the number of bits of source data, n denotes the number of bits of a codeword after modulation, End Zero (EZ) denotes the number of '0s' in a continuous sequence from the LSB of a codeword in a direction toward the MSB of the codeword, and LZ denotes the number of '0s' in a continuous sequence from the MSB of a codeword in a direction toward the LSB of the codeword. For example, codewords that satisfy d=0, k=10, m=8, n=15, $0\le EZ\le 8$ are divided according to the following LZ conditions:

- 1) number of codewords satisfying 2≤LZ≤10: 177
- 2) number of codewords satisfying 1≤LZ≤9: 257
- 3) number of codewords satisfying 0≤LZ≤6: 360
- 4) number of codewords satisfying 0≤LZ≤2: 262

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[0026] If the number of bits of source data satisfies m=8, the number of codewords for conversion should be 256 or more. However, in condition 1), the number of codewords does not amount to 256. Therefore, the number of codewords in condition 1) can amount to 256 by taking some codewords from a condition having surplus number of codewords. In this case, 83 codewords from the codewords satisfying group 3)'s LZ condition may be taken and added to group 1). Then, the numbers of codewords included in conditions 1) through 4) are 260, 257, 277(=360-83), and 262, respectively, and satisfy the minimum number of modulation codewords, that is, 256 for 8-bit source data. In the table of Figure 4, Main Code Group 1 (MCG1) is the name of a code group containing codewords satisfying condition 1) and some (83) codewords are taken from codewords satisfying condition 3). MCG2 and MCG4 are the names of codewords satisfying condition 2), and 4), respectively. MCG3 is the name of codewords satisfying condition 3), excluding the 83 codewords taken by MCG1. In each of the main code groups MCG1 through MCG4, only 256 codewords can be used as conversion codes for source codes.

[0027] Figure 5 is a table showing a variety of codeword groups of a conversion table for DC control and the characteristics of codewords in each codeword group. For example, assuming that d=2, k=10, m=8, n=17, and 0≤EZ≤8, conversion code tables for controlling DC components may include the following 4 groups (corresponding to DCG1, DCG2, DCG3, and DCG4 of Figure 5, respectively) according to the LZ conditions:

- 1) number of codewords satisfying 2<LZ<10: 375
- number of codewords satisfying 1≤LZ≤9: 546
- 1) number of codewords satisfying 0≤LZ≤6: 763
- 1) number of codewords satisfying 0≤LZ≤2: 556

[0028] Each group forming a conversion table for controlling DC components should have at least 2 codewords that selectively correspond to one source data, and therefore should have at least $512 (= 2^8 + 2^8)$ codewords for 8-bit source data. Since the number of codewords in the code group satisfying the LZ condition 1) is less than 512, code group 1) can take surplus codewords from other code groups satisfying other LZ conditions to amount to the number of 512. For example, in the above embodiment, code group 1) may take 177 codewords from the code group satisfying the condition 3) so as to have 552 (=375 + 177) codewords.

[0029] Figure 6 is a table showing a variety of codeword groups of an auxiliary conversion table for DC control and the characteristics of codewords in each code group. For example, among codewords satisfying d=2, k=10, m=8, and n=15, codewords satisfying 9≤EZ≤10, the remaining codewords of the main code conversion groups (MCGs), and codewords satisfying LZ=7, 8 or LZ=4, 5 are used as codewords of auxiliary code groups (ACGs) for controlling suppression of DC components. The conditions for generating these codewords will now be explained in detail. The following conditions correspond to ACG1 through ACG4, respectively, which are names of the auxiliary conversion tables for controlling suppression of DC components:

- 1) 5 codewords (satisfying $9 \le EZ \le 10$ and $LZ \ne 0$) + the remaining 4 codewords (in the MCG1) = 9 codewords,
- 2) 5 codewords (satisfying $9 \le EZ \le 10$ and $LZ \ne 0$) + 1 remaining codewords (in the MCG1) = 6 codewords,
- 3) 5 codewords (satisfying $9 \le EZ \le 10$ and $LZ \ne 1$) + 15 codewords (satisfying $7 \le LZ \le 8$ and $0 \le EZ \le 8$) = 41 codewords, the remaining 4 codewords in the MCG1 = 9 codewords,
- 4) 7 codewords (satisfying 9≤EZ≤10 + the remaining 6 codewords in the MCG4) + 85 codewords (satisfying 3≤LZ≤5 and 0≤EZ≤8) = 98 codewords.

[0030] When codeword a and codeword b are connected, the junction where the two codewords are connected should satisfy a run length (d, k) condition. Figure 7 is a diagram showing what should be considered for the run length conditions when codewords a and b are connected. Satisfying the run length condition means that in Figure 7 a value obtained by adding the end zero (EZ_a) of codeword a and the lead zero (LZ_b) of codeword b is equal to or greater than the minimum run length d and equal to or less than the maximum run length k.

[0031] Figure 8 is a table showing an example of changes in INV (whose meaning will be described below) before code conversion and after code conversion when the run length conditions described in Figure 7 are not satisfied. Codeword b is determined in a group indicated by the EZ of the preceding codeword, codeword a. When a or b is included in a code group which does not have enough codewords to meet the condition and takes codewords from other code groups, the (d, k) condition may not be satisfied. In this example, the EZ of codeword a changes to satisfy the run length condition, which is referred to as the boundary rule. Variable INV which indicates whether the number of '1s' in a codeword stream is an even number or an odd number may change from the previous INV while the codeword a didn't change (unclear), according to the boundary rule. Due to this characteristic, attention should be paid to allocation of a codeword between code conversion tables capable of controlling suppression of DC components.

[0032] Figure 9 is a diagram showing an example of code stream branching due to selective codewords b1 and b2 for DC control. One of the major features of the code conversion of the present invention is that the codewords of two code conversion tables that can be selected for DC control are allocated so that they have opposite INV characteristics. When the previous INV changes according to the boundary rule as described above, if the INVs of both codewords in the two code conversion tables that can be selected change, then there will be no problem. Otherwise, characteristics of codewords having opposite INV cannot be maintained. For this reason, a code conversion table is designed considering the following.

[0033] First, in A of Figure 9, that is, at the junction where the codeword a and the codeword b are connected to each other, if b1 and b2, which can be selected as codeword b, are codewords in DCG11 and DCG12, respectively, which are regrouped in the code conversion table DCG1 shown in Figure 5 after separating codewords which correspond to the same source code but have different INVs, or if b1 and b2 are codewords of MCG1 and MCG2, respectively, then codewords in which LZ_b1 (the number of LZs of codeword b1) and LZ_b2 (the number of LZs of codeword b2) is 1 are allocated on the location. By doing so, when the EZ of the codeword a is '0', according to the boundary rule, the INV of codeword a changes in both the code stream containing the codeword b1 and the code stream containing the codeword b2, or the INV of codeword a does not change in either the code stream containing the codeword b1 or the code stream containing the codeword b2, such that the INVs of the two code streams are maintained to be opposite. An example is as follows:

27 224 250 source data streaml(before conversion)000001000010001 (MCG3) 30 code 000001000001001 (MCG1) 010010010000000 (MCG1) conversion) 000001000010001 streaml(after code 35 010010010000000 000001000001000 0 1 INV1 stream2 (before conversion) 000001000010001 (MCG3) code 40 000001000001001 (MCG1) 01001000000000 (ACG1) conversion) 000001000010001 stream2(after code 010010000000000 000001000001000 45 1 1 1 INV2

[0034] Next, in B of Figure 9, that is, at the junction where codeword b and codeword c are connected to each other, if codewords b1 and b2 are respectively included in code conversion tables DCG11 and DCG12, DCG21 and DCG22, DCG31 and DCG32, DCG41 and DCG42, MCG1 and ACG1, MCG2 and ACG3, MCG3 and ACG3, or MCG4 and ACG4, and (xx)xxxxxxxxxx1001 or (xx)xxxxxxxxxx10001, INV may change according to the boundary rule due to the LZ of the following codeword c. Therefore, these codewords bland b2 are allocated to the location for corresponding same source data in each table such that the INVs of the two code streams are maintained to be opposite. An example is as follows:

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	source	e data	250	. 152		210	
5	code s	stream1(before	conv	ersion)	0000010	00010001 (MC	:G3)
_	010000	000010001001 (DC	CG11)	000001000	00001 (MCG	1)	
	code	stream1(after	conv	version)	0	00001000010	000
10	010000	000010001001		000001000	00001		
	INV1		0	0		0	
	code s	stream2(before	conve	ersion)	0000010	00010001 (MC	:G3)
15	010010	000010001001 (DC	G12)	0100000010	01001 (MCG	1)	
	codes	tream2(after	conv	rersion)	00	00010000100	000
	010010	000010001001		0100000010	01001		
20	INV2		0	1		1	

[0035] For the junctions A and B of Figure 9, the codewords are first allocated to the location corresponding same source data in each code conversion table (DCG11 and DCG12 or MCG1 and ACG1) considering above. Referring to the following example, in point B, according to the boundary rule, the INVs of code stream1 and code stream2 are maintained to be opposite and the INVs of code stream3 and code stream4 are maintained to be opposite. Also, at point B, according to the boundary rule, the INVs of code stream1 and code stream3 are maintained to be opposite and the INVs of code stream2 and code stream4 are maintained to be opposite.

source data 250 152 7 code stream1(before conversion) 000001000010001 (MCG3) 35 0100000010001001(DCG11) 010000010010001(MCG1) code stream1(after conversion) 000001000010000 0100000010001000 010000010010001 40 INV1 0 1 1

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	code	stream2(before	conve	ersion)	0000	01000010001 (MCG3)
5	01000	000010001001 (DC	CG11)	0100100100	10001(<i>F</i>	ACG1)
	code	stream2(after	conv	rersion)		0000010000100000
	01000	000010001000		0100100100	10001	
10	INV2		0		1	0
	code	stream3(before	conve	ersion)	0000	01000010001 (MCG3)
	01001	.000010001001 (DC	CG12)	0100000100	10001 (N	(CG1)
15	code	stream3(after	conv	version)		000001000010000
	01001	.000010001000		0100000100	10001	
	INV3		0		0	0
20	code	stream4 (before	conve	ersion)	0000	01000010001 (MCG3)
	01001	.000010001001 (DC	CG12)	0100100100	10001(2	ACG1)
	code	stream4(after	conv	rersion)		0000010000100000
25	01001	.000010001000		0100100100	10001	
	INV4		0		0	1

30 [0036] As described above, considering changes in the INV of a codeword due to the boundary rule in a codeword stream, codewords are allocated so that the INV polarities of a codeword pair after modulation is always be maintained to be opposite. Figure 10 is a graph showing the relationship of INV values of this code stream pair. If codewords are allocated such that the INV values of a code stream pair are always opposite, a codeword can be selected so that a code stream which is DC components between the code stream pair is formed.

[0037] Exceptions to the rule that INV values are maintained to be opposite at point A of Figure 9 may occur when source data is from 251 to 255 in the code conversion table for controlling DC components. In such exceptional cases, the CSV signs of codewords are made to be opposite so that the difference between DSV values in the code stream pair is made.

[0038] Figures 11a through 11e are main conversion code tables in which codewords are generated and allocated considering conditions described above.

[0039] Figures 12a through 12j are code conversion tables for DC control in which codewords are generated and allocated considering conditions described above.

[0040] Figures 13a and 13b are auxiliary code conversion tables for DC control in which codewords are generated and allocated considering conditions described above.

[0041] Figure 14 is a graph showing the difference between the frequency spectrum when codewords of the code conversion table for DC control according to the present invention are used in 25% of all of the codewords, and the frequency spectrum when prior art EFM+ modulation codewords are used. The graph shows that in a low frequency band, the frequency spectrum of the modulated code stream according to the present invention is almost the same as the frequency spectrum of the EFM+, which indicates that the code stream of the present invention has almost the same capability of suppressing DC components as that of the EFM+ method.

[0042] Accordingly, since the present invention uses 15-bit codes as the main conversion code and selectively uses 17-bit DC control codes for controlling DC components, the present invention has better efficiency in recording density than the prior art EFM+ code and has the same DC suppression capability as the EFM+ code.

[0043] In the present invention, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density.

[0044] Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

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[0045] The reader's attention is directed to all papers and documents which are filed concurrently with or previous to this specification in connection with this application and which are open to public inspection with this specification, and the contents of all such papers and documents are incorporated herein by reference.

[0046] All of the features disclosed in this specification (including any accompanying claims, abstract and drawings), and/or all of the steps of any method or process so disclosed, may be combined in any combination, except combinations where at least some of such features and/or steps are mutually exclusive.

[0047] Each feature disclosed in this specification (including any accompanying claims, abstract and drawings), may be replaced by alternative features serving the same, equivalent or similar purpose, unless expressly stated otherwise. Thus, unless expressly stated otherwise, each feature disclosed is one example only of a generic series of equivalent or similar features.

[0048] The invention is not restricted to the details of the foregoing embodiment(s). The invention extends to any novel one, or any novel combination, of the features disclosed in this specification (including any accompanying claims, abstract and drawings), or to any novel one, or any novel combination, of the steps of any method or process so disclosed.

Claims

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 A method for generating and allocating codewords of source words which are to be recorded on a recording medium, the method comprising:

generating codewords satisfying predetermined run length conditions and grouping codewords according to each run length condition; and

allocating the codewords such that a codeword for the source word is capable of controlling suppression of DC components.

2. The method of claim 1, further comprising:

when a predetermined boundary condition is not satisfied in the code stream, allocating codewords such that codewords which satisfy the boundary condition and maintain the DC control characteristics which are considered when the initial codewords are allocated replace the initial codewords.

- 3. The method of claim 1 or 2, wherein in order to make code streams capable of controlling suppression of DC components allocating codewords such that a pair of codeword streams having opposite INV characteristics are made selectable, where INV indicates whether the number of '1s' is an odd number or an even number.
- 4. The method of claim 1, 2 or 3, wherein the step for generating codewords comprises:

generating codewords satisfying the length of a predetermined first codeword, and predetermined run length conditions, grouping the codewords according to each predetermined run length condition to generate a main conversion codeword table;

generating DC control codewords satisfying the length of a predetermined second codeword, and predetermined run length conditions in order to control DC components in the code(word) stream, grouping the DC control codewords, and to generate a code conversion table for controlling DC components; and

generating additional DC control codewords by taking codewords which satisfy the predetermined run length conditions and are not needed in the main conversion codeword table, and grouping the additional DC control codewords.

- 5. The method of claim 4, wherein when the bit length of the source word is 8, the length of a codeword in the main conversion code table is 15 bits.
- 6. The method of claim 5, wherein the main conversion code table contains groups of codewords, the groups formed of a group of codewords each having from 2 to 10 Lead Zeros (LZs), a group of codewords each having from 1 to 9 LZs, a group of codewords each having from 1 to 6 LZs, and a group of codewords each having from 0 to 2 LZs, while having from 0 to 8 End Zeros (EZs).

- 7. The method of claim 6, wherein among the groups of the main conversion code table, a group having a lesser number of codewords than the minimum number of codewords for converting the source data takes surplus codewords from a group having a greater number of codewords than the minimum number of codewords so as to amount to the minimum number of codewords.
- 8. The method of claim 5, wherein the length of the codewords of the DC control conversion code table is 17 bits.
- 9. The method of claim 8, wherein the DC control conversion code table contains groups of codewords, the groups formed of a group of codewords of which LZ is from 2 and to 10, a group of codewords of which LZ is from 1 to 9, a group of codewords of which LZ is from 0 to 6, and a group of codewords of which LZ is from 0 to 2, while EZ is from 0 to 8.
- 10. The method of claim 9, wherein each code group of the DC control conversion table has as much codewords as a source word can correspond to pairs of codewords, each pair of codewords has opposite INV characteristics and is selectable.
- 11. The method of claim 10, wherein among the groups of the DC control conversion code table, a group having a lesser number of codewords than the minimum number of needed codewords takes surplus codewords from a group having a greater number of codewords than the minimum number of codewords, so as to have equal to or greater than the minimum number of codewords.
- 12. The method of claim 8, wherein the auxiliary DC control conversion table having a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having at least one LZ, and codewords which are taken from surplus codewords of the first main conversion code group;
 - a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having at least one LZ, and codewords which are taken from surplus codewords of the second main conversion code group;
 - a group of codewords, each having a length of 15 bits, having from 9 to 10 EZs, and having no LZ, the surplus codewords of the third main conversion code group, and having 7 to 8 LZs, or having 0 to 8 EZs;
 - a group of codewords, each having a length of 15 bits, and having 9 or 10 EZs, the surplus codewords of the fourth main conversion code group, and having 3 to 8 LZs and 0 to 8 EZs.
- 13. The method of claim 12, wherein when the code stream pairs are a, b1, c and a, b2, c, respectively, and b1 and b2 are DC control codewords having opposite INV characteristics, codewords are allocated such that the INV characteristics of code streams after conversion are maintained to be opposite even if code changes of a, b1, b2, or c occur due to violation of the predetermined run length between a, and b1(b2) or b1(b2) and c.
- 40 14. An allocation method for allocating codewords generated for source words to be recording on a recording medium, the method comprising:
 - when a preceding codeword a and a following codeword b form a code stream X, allocating one of two selectable codewords b1 and b2 as codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number; and
 - when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords.
 - 15. The method of claim 14, wherein when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword a in the Most Significant Bit (MSB) direction is 0, and the number of continuous '0s' from the MSB of the codewords b1 or b2 in the LSB direction is 1, code changes of either the codeword a or b1(b2) occur to satisfy the boundary condition.
 - 16. The method of claim 14, wherein when the number of continuous '0s' between the codewords a and b is 1 or 0,

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the codeword a or b is changed such that the number of 0s forming the boundary is greater than 2 and less than 10.

- 17. The method of claim 16, wherein the codeword a of the code stream X1 and the codeword a of the code stream X2 are changed to other codewords such that the resulting codewords a of code streams X1 and X2 have the same INV value, and as a result, by the INVs of codewords b1 and b2 following the codewords a respectively, the INVs of the X1 and X2 become different.
- **18.** An allocation method for allocating codewords of source words to be recording on a recording medium, the method comprising:

when a preceding codeword b and a following codeword c form a code stream Y, allocating one of two selectable codewords b1 and b2 as the codeword b, wherein codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number; and

when the code stream of b1 and c is Y1, and the code stream of b2 and c is Y2, allocating codewords such that INVs of Y1 and Y2 are maintained to be opposite when the codeword b1, b2 or c should be replaced by another codeword in compliance with a predetermined boundary condition between codewords.

- 19. The method of claim 18, wherein when the predetermined boundary condition is that the number of continuous '0s' between codewords should be at least 2, and when the number of continuous '0s' from the Least Significant Bit (LSB) of the codeword c toward the Most Significant Bit (MSB) is 1, codeword b which does not satisfy the boundary condition and is xxxxxxxxxxxx1001 or xxxxxxxxxx10001 appears in both b1 and b2.
- 25 20. The method of claim 18, wherein when the number of continuous '0s' between the codewords a and b is 1 or 0, the codeword a or b is changed such that the number 0s forming the boundary is greater than 2 and less than 10.

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FIG. 1

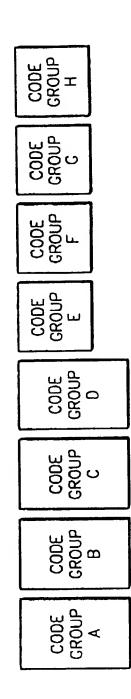


FIG. 2

NCC		_		2		2		4
CODE GROUP CONTAINING CODEWORD INDICATED BY	CONVERSION CODE GROUP	OC SUPPRESSION CODE GROUP FOR OC SUPPRESSION CONTROL WITH THE HELP OF MCCI	CONVERSION CODE GROUP	OC SUPPRESSION CODE GROUP FOR OC SUPPRESSION CONTROL WITH THE HELP OF MCG2	CONVERSION CODE GROUP	DC SUPPRESSION CODE GROUP FOR DC SUPPRESSION CONTROL WITH THE HELP OF DCC1	CONVERSION CODE GROUP	CONVERSION CODE GROUP FOR CONVERSION CODE GROUP FOR CODE CONTROL WITH THE GROUP CONTROL WITH THE GROUP GROUP FOR CONTROL WITH THE GROUP GR
NCC	190W	1st DSV CODE GROUP	MCG2	2nd DSV CODE GROUP	1920	3rd DSV CODE GROUP	2900	4th DSV CODE GROUP
CHARACTERISTIC LZ=2~9	6~2=77	6~2=71	1~0=71	1~0=Z	b15 (MSB)=b3=0)=63=0	b15 (MSB):	b15 (MSB)=1 EE b3=1
METHOD FOR DUPLICATED CODE		W3dDO	ROD OF EZ=2	~5 REPEATEDLY OC	CURS IN ALL K	CODEWROD OF EZ=2~5 REPEATEDLY OCCURS IN ALL KINDS OF CODE GROUPS	uPs	

FIG. 3

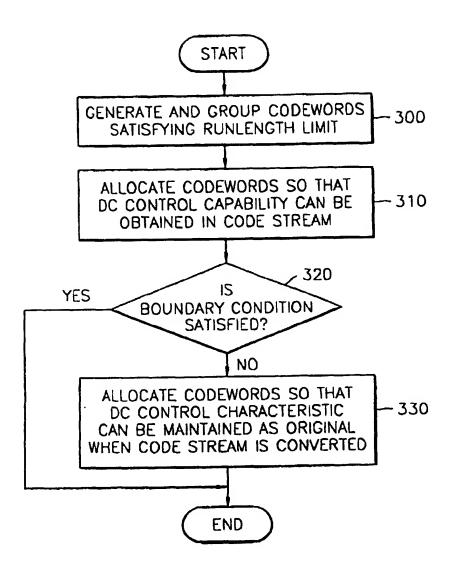


FIG. 4

	LZ (EZ)	NUMBER	ADD	NUMBER	DELETE	NUMBER	TOTAL NUMBER (DUPLICATED)
MCG1	2≤LZ ≤10 (0≤EZ≤8)	177	LZ=1 (0≤EZ <u>≤</u> 8)	83			260(4)
MCG2	1 <u>≤</u> LZ <u>≤</u> 9 (0≤EZ <u>≤</u> 8)	257	-		·	·	257(1)
MCG3	0≤LZ≤6 (0≤EZ≤8)	360			LZ=1 (0≤EZ≤8)	83	277(21)
MCG4	0≤LZ≤2 (0≤EZ≤8)	262					262(6)

FIG. 5

\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\\	LZ . (EZ)	NUMBER	ADD	NUMBER	DELETE	NUMBER	TOTAL NUMBER
DCG1	2≤LZ≤10 (0≤EZ≤8)	375	LZ=1 (0≤EZ≤8)	177			552
DCG2	1≤LZ≤9 (0≤EZ≤8)	546					546
DCG3	0≤LZ≤6 (0≤EZ≤B)	763			LZ=1 (0≤EZ≤8)	177	586
DCG4	0≤LZ≤2 (0≤EZ≤8)	556					556

FIG. 6

	LZ (EZ)	NUMBER	ADD	NUMBER		NUMBER	TOTAL NUMBER
ACG1	LZ≠0 (9≤EZ≤10)	5	SURPLUS CODE OF MCG1	4			9
ACG2	LZ≠0 (9≤EZ≤10)	5	SURPLUS CODE OF MCG2	1			6
ACG3	LZ≠1 (9≤EZ≤10)	5	SURPLUS CODE OF MCG3	21	7≤LZ≤8 (0≤EZ≤8)	15	41
ACG4	(9≤EZ≤10)	7	SURPLUS CODE OF MCG4	6	3≤LZ≤5 (0≤EZ≤8)	85	98

FIG. 7

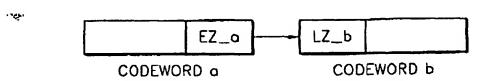


FIG. 8

CODEWORD a	CODEWORD 6	CHANGE IN INV
xxxxxxxxx001001 (BEFORE CONVERSION) xxxxxxxxx001000 (AFTER CONVERSION)		
xxxxxxxxx010001 (BEFORE CONVERSION) xxxxxxxxx010000 (AFTER CONVERSION)		CHANGE
xxxxxxxxx100001 (BEFORE CONVERSION) xxxxxxxxx100100 (AFTER CONVERSION)	0100xxxxxxxxxx	
~		NO
xxx10000000001 (BEFORE CONVERSION) xxx100000000100 (AFTER CONVERSION)		

FIG. 9

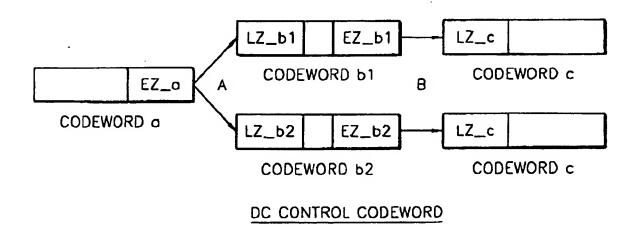


FIG. 10

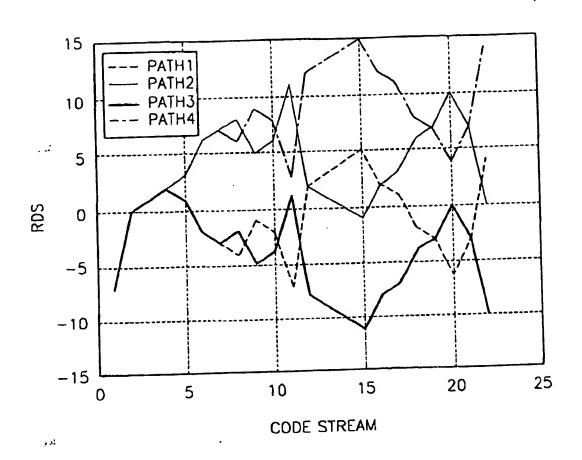


FIG. 11A

DATA	NCG1		NCC:		MCG3		NCG4	
SYMBOL	Code Word MSB LSB	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code Word MSB LSB	NCG
000	000100100000000	4	000100100000000	4	1001000000000001	1	100100000000001	1
000	000100100000000	4	000010010000000	1 4	0010010000000001	j	0100100000000001	li
001 002	0010001001000000	4	000100100100000	4	000000100000001	l i l	0010010000000001	li
002	010001001000000	4	00000010000000	4	100000000100001	l i l	100100100000001	l i
003	010000100100000	4	000000010000000	4	100000001000001	i	100010010000001	l i
005	010000001001001	i	000000000100010	2	1001000000000010		100001001000001	. 1
006	010000010010000	3	0001000000000001	l ĭ ˈ	0010010000000010	2 2 2 2	100000100100001	1
007	010000010010001	1	010010010000001	1	0010000000000010	2	100100010000001	1
008	010010000010010	2	0100010010000001	1	100100000010010	2	100010001000001	1
009	010010010000001	1	010000100100001	1 1	001001000010010	2	100001000100001	1
010	000001001000000	4	0000100000000001	1	000100100010010	2 2	0100100100000001	[]
011	000000100100000	4	010010001000001	1	1000000000010010	2	010001001000001	1 1
012	000000010010000	3	010001000100001	1	100000000100010	23333	010000100100001	1 1
013	000000001001000	3	001001001000001	1 1	100100000000100	3	100000000100001	
014	000000000100100	3	001000100100001	1	001001000000100	3	1001000000000010	1 6
015	000000000100010	2	0000010000000001		000100100000100	3	00100000000000010	1 2
016	000000000100001	3	001001000100001	1 1	000000100010000	3	1001001000000010	1 5
017 018	000000000010000	i	000100100100001	1	10000001001000	3	1000100100000010	1 5
019	010010001000000	4	000000100000001	l í	100000000100100	3	100001001000010	2
020	010010000000010	2	0100100000000001	l i i	100000100010000	3	100000100100010	1 2
021	010000100100001	ī	0010010000000001	ī	100000010001000	13	100000010010010	2
022	010001000100000	4	0001001000000001	1	100000001000100	3	100100010000010	2
023	010010000000001	1	0000100100000001	1	100001000010000	3	100010001000010	2
024	000100000000001	1	000001001000001	1	100000100001000	3	100001000100010	2
025	001001001000000	4	000000100100001	1	100000000100000	4	100000100010010	2
026	000000000010001	1 1	010001000000001	1	000100100000000	4	010010010000010	1 5
027	010010010000000	4	001000100000001 0001000100000001	1	000010010000000	4	1001000000100010	1 2
028 029	0000100000000001	3	000100010000001	l i	00000100100000	4	010010000010010	12222222222222222223
030	010010001000001	ĭ	000001000100001	i	100010010000000	4	100000000010010	2
031	010000000100100	3	0100001000000001	l i l	100001001000000	4	100000000100010	2
032	010001000100001	1	0010000100000001	i	100000100100000	4	010000000010010	2
033	010000100010000	3	0001000010000001	1	100100010000000	4	100000001000010	2
034	010000000010010	2	000010000100001	1	100100000001001	1	010000000100010	1 2
035	010000000001001	1	010000010000001	1	001001000001001	1 1	100100000000100	3.
036 <u>3</u> 037	001001000000001	1 4	001000001000001 000100000100001	1	000100100001001 100000010010001	1 1	0100000000000100	3
037	001000100000000	2	010000000100001	l î	100000001001001	l î '	100000000010000	3
039	********	ī	010000001000001	l î	100000100010001	l i	010000000001000	3
040	000100010000000	4	001000000100001	l i	100000010001001	l ī	001000000000100	3
041	010000100010001	1	000000000100001	1	100010010000001	1	100100100000100	3
042	000010010000001	1	000000001000001	1	100001001000001	1	100010010000100	1 3
043	000000100010000	3	000000010000001	1 1	100000100100001	1	100001001000100	33333333333
044	010010000100000	4	0010000000000010	2 2	0001000000000001	1 1	100000100100100	1 3
045	000001001000001	3	0001000000000010	2 2	100010001000001	1 1	010000000010000	1 3
046	000000010001000	1	010010010000010	2	0000100000000001	li	100100100001000	1 3
047 048	000000000000000000000000000000000000000	3	010000100100010	2	100100001000001	l i l	100100010000100	3
049	000000001000100	2	010000010010010	2	100010000100001	l i l	100010010001000	3
050	010010000001001	ī	000010000000010	2	001001001000001	l i l	100010001000100	1 3
051	0100100000000100	3	010010001000010	2	001000100100001	i	100100010010000	3
052	000010001000000	4	010001000100010	2 2	0000010000000001	1	100100001001000] 3
053	0001000000000010	2	010000100010010	2	100100000100001	1	100100000100100	3
054	000001000100000	4	001001001000010	2	001001000100001	1	010010010010000	1 3
055	010000010001001	1	001000100100010	2	000100100100001	1	100000010010000	3
056	000001000010000	3	001000010010010	2	000100100000001		100000001001000	3
057	001001000100000	4	000001000000010	2	000010010000001]]	100000000100100	3
058	000000100001000	3	010010000000010	2	000001001000001	1 1	100000100010000	333333333
059	000000010000100	9 2	001001000000010 0001001000000010	2 2	1000100000000001	li	10000001000100	3
060	000000000010010	6	OTOMOTOTOM	-	TOTAL		10000001000100	

FIG. 11B

	NCG1		NCG2		MCG3		NCG4	 -
DATA SYMBOL	Code Word	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code Word	NCG
	MSB LSB 001001000001001	1	MSB LSB 000010010000010	2	001000100000001	1	010000010010000	3
061 062	0010010000001001	3	000001001000010	2	000100010000001	1 1	010000001001000	3
063	010010010000010	2	000000100100010	1 2 1	000010001000001	j	010000000100100	3
064	001001001000001	1	000000010010010	2	000001000100001	1	001000100000000	4
065	000100100100000	4	010001000000010	2	1000010000000001	i	001000010000000	4
066	001001000000010	2	001000100000010	2	000100001000001	lil	100100100000000	4
067	010001001000010 001000100100001	2	000010001000010	2 2	000010000100001	lil	1000100100000000	4
860 069	000000100100000	4	00001000100010	2	100000100000001	1	100001001000000	4
070	000000010010001	li	000000100010010	2	001000001000001	1 1	100000100100000	4
071	010000100100010	2	010000100000010	2 2 2	000100000100001	1 1	100100010000000	4
072	001000010010001	1	001000010000010	2	100000010000001		100001000100000	1 4
073	010000010001000	3 2	000100001000010	2	100010010000010	2	010010010000000	4
074	010000010010010	1	000010000100010	2	100001001000010	2	010001001000000	4
075 076	010000001000100	3	010000000010010	2	100000100100010	2 2 2	010000100100000	4
077	0000100000000010	2	010000000100010	2	100000010010010	2	100100001000000	4
078	000001000000001	1	001000000010010	2	000100000000010	2 2	010010001000000	4
079	001000010010000	3	010000001000010	2 2	10001000100010	2	010001000100000	4
080	010000000100010	2	001000000100010	2	100000100010010	2	001001001000000	4
081 082	010000000000000000000000000000000000000	3	010000010000010	2	000010000000010	2 2	100100000001001	1
083	000100100000010	2	001000001000010	2	100100001000010	2	100000000001001 100000000010001	1 1
084	001000001001000	3	000100000100010	5	100010000100010	2 2	010000000000000000000000000000000000000	li
085	000000001001001	1	000010000010010	2 2	100001000010010 001001001001001001001001	2	010000000010001	i
086	001000000100100	3	0100100001000010	2	001000100100010	2	001000000001001	1
087 088	010000100000000	4 2	001001000010010	2	001000010010010	2	100100100001001	1 ;
089	001000000001001	lī	000100100010010	2	000001000000010	2	100010010001001	1
090	{ 01000000001000	3 2	000010010010010	2	000100100000010	2 2	001000000010001	l i
091	000010010000010		000000010000010	2 2	000001001000010	2	100000010010001	1
092	0100010000000001	3	010010000100010	2	000000100100010	2	100000001001001	1 1
093 094	010010001001000	2	001001000100010	2 2	1000100000000010	2	100000100010001	1 1
095	000001001000010	2	001000100010010	2	001000100000010	2	100000010001001	i
096	1 0010001000000001	1	000100100100010	2 2 3	00010001000010	2 2	010000001001001	li
097	010010000100100	3	000100010010010	2	000001000100010	2	100100001000001	1
098 099 *	010010000100001	1 2	010000000000100	1 3	000000100010010	2 2	100010000100001	1
100	000100010000001	li	010000000001000	13	100001000000010	2	010010001000001	1
101	001001001001000	3	001000000000100	3	001000010000010	2	010001000100001	li
102	000000010000000	4	010000000010000	3	000100001000010	2 2	001000100100001	Ī
103	001001000010010	2	00100000001000	3	000001000010010	2	100100000100001	1
104	0001001000001001	3	010010010000100	3	100000001000010	2	010010000100001	1
105 106	010001000010000	l ă	010001001000100	1 3	001000000010010	2	001001000100001	1
107	000010001000001	1	010000100100100	3	100010010010010	2 2	100010000000001 0100010000000001	l i
108	010000100001000	3	001000000010000	3	100000010000010		001000100000001	li
109	001000010000000	4	000100000001000	3	0001000000100010		100001000000001	1
110	010000001000010	2	000010000000100		100000100000010		010000100000001	1 1
111	010000000100001	3	001001000000100	3	001000001000010	2	001000010000001	11
112 113	000001000100001	li	1 000100100000100	3	000100000100010	2	100000100000001	1 1
113	001001000100100	3	000010010000100	3	000010000010010	2	010000010000001	li
115	010001000100010	2	000001001000100	3	000010010010010	2 2	10000001000001	
116	000000100010001	1 1	000000100100100	3	10010000010010		1 010000000100001	1
117	000100100100100	3	010010000001000 010001000000100	3	001001000100010	2	100000010000001)
118 119	010001000010001	1	1 001001000001000	3	001000100010010) 2	010000001000001	1 1
120	000010000010000	3	001000100000100	3		2	001000000100001	
160	COCCETOCOCCETOCOC							

FIG. 11C

DATA	MCG1		JICG2		NCG3		MCG4	
SYMBOL	Code Word	NCG	Code Word	NCG	Code Word	NCG	Code Word	NCG
-	MSB LSB	-	NSD LSD	-	MSB LSB		MSB LSB	
121 122	000000000100000	4 2	000100100001000	3	000100010010010	2	010000100100010	2
123	00010010001001	li	000100010000100	3	000000100000010 100000000001000	2	010000010010010 100100001000010	2 2
124	000010010000100	i	000010001000100	3	10000000001000	3	10010000100010	5
125	01000010000100	3	000001001001000	3	001000000000100	1 3	100001000010010	2
126	000100001000000	4	000001000100100	3	100010010000100	3333333	010010001000010	222222222222222222222222222222222222222
127	001000000100010	2	01000001001(000	3	100001001000100	3	010001000100010	2
128	001000000010001	1	010000001001000	3	100000100100100	3	010000100010010	2
129 130	010000000010000	3	010000000100100	3	001000000001000	3	001001001000010	2
131	010000100010001	2	010000010001000	3	000100000000100 100010010001000	3	001000100100010 001000010010010	1 5
132	000000100000100	3	010000001000100	3	1000100010001000	3	0010010000000010	1 5
133	010000100001001	ì	001000010010000	3	100001001001000	3	100010000000010	2
134	010010010001000	3	001000001001000	3	100001000100100	3	0100010000000010	l ž
135	000000001000000	4	001000000100100	3 3	001000000010000	.3	001000100000010	2
136	000010010010010	2	010001000010000] 3	000100000001000	3	100001000000010	2
137 138	000001001001001	3	010000100001000	3 3	000010000000100	3	010000100000010	1 2
139	000100000010010	2	001000100010000	3	000010010000100 000001001000100	ოოოოო	001000010000010	5
140	0001000000001001	ĩ	001000010001000	3	00000100100100	3	100010010010010	2
141	001000000001000	3	001000001000100	3	100100000001000	3 1	100000010000010	Ž
142	001001001000010	2	000100010010000	3	100010000000100	I 3 I	010000001000010	2
143	001001000100001	1	000100001001000	3	001001000001000	3	001000000100010	2
144 145	001000100010000 000000010010010	3 2	000100000100100	3	001000100000100	3	100100010010010	2
146	010001001001001	ĺ	010001000001000	3	000100100001000 000100010000100	3	010010010010010	5
147	0001000000000100	3	010000100000100	3	000010010001000	3	010000010000010	ž
148	001000100100010	2	001001000010000	3	000010001000100	3	001000001000010	2
149	001000100010001	1	001000100001000	3	000001001001000	3	001001000010010	2
150 151	001000010001000	3	001000010000100	3	000001000100100	3	100100000100010	2
152	0100001000000001	1 3	000100100010000 000100010001000	3	100000010000100	3	100010000010010	2
153	001000010010010	2	000100001000100	3	001000001001000	3	0100010000100010	2
154	001000010001001	1	000010010010000	3	001000000100100	3	001001000100010	2
155	001000001000100	3	000010001001000	3	100010000010000	3 3	001000100010010	2
156	000010000100000	4	000010000100100	3	100001000001000	3	100001001001000	3
157 158	010000010000010 010000001000001	2	000000010010000	3	100000100000100	3	100001000100100	3
	010000000100000	4	0000000000100100	3	001000010001000	3	010001001000100	จั
160	000100010010000	3	000000100010000	ă l	001000001000100	3 3 3	010000100100100	3
161	000000100100100	3	000000010001000	Š	000100010010000	3	001000000010000	3
162	000100100100001	1	000000001000100	3	000100001001000	3	010010000000100	3
163 164	000100001001000	3 2	000001000010000	3	000100000100100	3	0010010000000100	3
165	0010000010000010	2	000000100001000	33333333	100100000010000	3	100100000001000	3
166	001000010000001	ī	010010010010000	3	10001000000100	3	010010000001000	3
167	001000000100001	ī	010010001001000	3	001001000010000	3	010001000000100	3
168	010010001000100	3	010010000100100	3	001000100001000	3	001001000001000	3
169	010001001000100	3	001001001001000	3	001000010000100	3	001000100000100	3
170	000000001000001	1	001001000100100	3	000100100010000	3	100001000010000	3
171	000100000100100	3	000100100100100	3	000100010001000	3	100000100001000	3
172 173	010010000010000	3	000010000010000	3	000100001000100	3	010000010001000100	ğ
174	001000100000010	ž	00000100000100	š	00001001001000	3	010000010001000	3
175	000100000100010	2	010010010001000	9	000010000100100	3	010000001000100	3 3 3 3
176	000100001000001	ī	010010001000100	3	000001000010000	3	001000010010000	3
177	000100000010001	1	010001001001000	3	000000100001000	3	001000001001000	3
178	010001001001000	3	010001000100100	3	100100001001000	3	001000000100100	3
179 180	010000100100100	3	001001001000100	3	100100000100100	9	100010000010000	3
100	ATMATAMAMATON		201000100100	- 1	001001001001		***************************************	 _

FIG. 11D

SPANEON		MCG1		MCC2		NCG3		MCG4	
No. No. No		Code Word	NCG		NCG		NCG		NCG
182		11-10-	2	1400	7	III.D.D	3	100000100000100	
1000 1000	181	00010001000010		00010000001000			3	010001000010000	
185		0100010000100001		000001000000100	3		3]
188		000001000000010	2				3		3
188		000000010000010	2		4	000000100000100	3		1 3
188				000000001000000	4	100100001000100	3		3
214		010010000010001	1			100010010010000	3		3
214		010001000001000					3		3
214			2	010000100100000			3	100010000001000	١ă
214							3	010010000010000	3
214							ă		3
214			2				3	010000100000100	3
214		000010000010010				000001000000100	3	001001000010000	3
214						100000001000000	4	001000100001000	3
214				000001001000000		100010001000000			3
214				000000100100000		100001000100000			3
214		001001000010000	3	001000100000000		100100001000000			1 3
214		000000100010010				100010000100000			3
214						001001001000000			3
214								100100010001000	l š
214		010010001001001			1				3
214			1 3						3
214		00100000001000	3	000010000100000		000000100000000	4	100010001001000	1 3
214		01000010000010	l î	010000010000000		001000100000000		100010000100100	3
214		000100000010000		001000001000000		000100010000000		010010010001000	1 3
214		010010000100010	1 2			000010001000000		010010001000100	1 3
214		010001000010010	2					01000100100100100	3
214	210	000000100000001				00100001000000		001001001000100	3
214		010001000001001				000100001000000		001000100100100	3
214		001000100001000	3			10000100001000000		100000000100000	4
215		00100100000100	1 3		l î	001000001000000	4	100000001000000	
216				001000010010001		000100000100000	4	010000000100000	
217					1				4
218		000010000001000	1 3	010001000010001	1			100100000100000	
219		000010000000100 .	3		1			011011000100000	
221					1	1000001001001			
222								010000010000000	4
223 00100010000100 1 0100000001001 1					1 1			001000001000000	
224 00000100000000							1		
225 01000001000000						001000100010001			
226		01000010000000				001000010001001		001000000100000	
227				001000000010001			-	100001000010001	
228 001001000010001 1 01001001001001 1 1000000001001 1 0100001001001 1 0100000001001 1 0100000001001 1 010000001001001 1 00100000001001 1 00100000001001 1 0010000001001001 1 0010000001001 1 0010000001001 1 0010000001001 1 0010000001001 1 0010000001001 1 0010000001001 1 0010000001001 1 100001000001001 1 100001000001001 1 100001000001001 1 100001000001001 1 100001000001001 1 100001000001001 1 100001000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000001001 1 1000010000000001 1 1000010000000001 1 1000010000000000000000000 1 1000010000000000000000000000000000000			2			000100001001001			
229 001000100001001 1 010001001001001 1 00100000010001 1 00100000010001 1 0010000001001001 1 0010000001001001 1 0010000001001001 1 0010000001001001 1 001000001001001 1 001000001001001 1 001000001001001 1 100001001001001 1 100001001001001 1 100001001001001 1 100001000001001 1 100001000001001 1 1000010000001001 1 1000010000001001 1 10000100000010001 1 10000100000010001 1 10000100000010001 1 10000100000010001 1 10000100000000001 1 100001000000010001 1 10000100000000001 1 100001000000000001 1 10000100000000001 1 10001000000000001 1 100010000000000001 1 100010000000000000000001 1 1000100000000000000000000000000000000		001001000010001	1	010010010001001		10000000000001001			
231	229	001000100001001	1	010001001001001		100000000000000000000000000000000000000		0010000010010001	
231 00100010001000 3 0100100001001 1 1000100100101 1 1000100001001	230	000100100010000				1000100000001001		001000001001001	1
233		001000100000100				10001001001001		100010000010001	1
234				010010000001001		001001001001001		100001000001001] 1
234	233	001000001000000		000100100001001				010001000010001	
235		000100001000100		00010010001001				1 010000100001001	
237 000010001001000 3 01001000010001 1 00010000010001				000001001001001		100010001001001	1	001000100010001	1 1
238 01000001000000 4 0100100001001 1 00001000001001 1 10001001				010010000010001		000100000010001	1	001000010001001	1 1
239 00010010010010 2 00100100001001 1 00001001001 1 10001001		010000001000000		010001000001001		000010000001001		100100010001001	1 1
		000100100100010	2	001001000010001		000010010001001		100010001001001	li
240 000100010010010 2 00100010000100	240	000100010010010	2	001000100001001	11	000001001001001		10001001001001	

FIG. 11E

DATA	MCG1		MCG2		MCG3		MCG4	
SYMBOL	Code Word	NCG	Code Word MSB LSB	NCG	Code Word MSB LSB	NCC	Code Word MSB LSB	NCG
241 242 243 244 245 246 247 248 249 250	000100100010001 000100010001001 00001000010001000 0001001	113341213233	000100100010001 000100100010001 00001001	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	100100000010001 100010000001001 00100100	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	010010010001001 010001001001001 01001000000	1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1
253 254 255	0000100100100 000000010000001 0000010001001	3 1 3	001001001001001 000010000010001 000001000001001	1 1 1	100100001001001 001001001001001 000010000010001 000001000001001	1 1 1	100100001001001 01001001001001 010010001001	1 1 1 1 1

FIG. 12A

2471	DCG11		DCG12		DCG21		DCG22	
DATA SYMBOL	Code Word	NCG	Code Word MSB 1.SB	NCG	Code Word NSB LSB	NCG	Code Ford NSB LSB	NCG
	MSB LSE	<u>'</u>	001001000000000001	1	01001001000000001	1	001001000000000001	3
000	0010010010000000	li	0001001000000000001	l i	01000100100000001	1	000100100000000001	1 1
001	0010000100100000	li	000010010000000001	Ιi	101000010010000001	1	00001001000000001	1 1
002 003	0010000010010000		000001001000000001	1	01000001001000001	1	00000100100000001	1 ;
004	0000010000000000		00000010010000001)	01000000100100001	1	00000010010000001	1:
005	0010010001000000		00000001001000001	1	010010001000000001	1	00000001001000001	1 1
006	0010001000100000	1	00000000100100001	1 1	01000100010000001		01000100000000001	l i
007	0010000100010000	1 1	001000100000000001	1 ;	01000010001000001	1	00100010000000001	1 1
800	0001001001000000	1	000100010000000001	1 1	001001001000000001	١i	000100010000000001	1
009	0001000100100000]]	000010001000000001		00100010010000001		1 0000 1000 100000001	1
010	0001000010010000	! !	00000100010000001		00100001001000001	li	00000100010000001	1
011	0000001000000000		00000001000100001		00100000100100001		100000010001000001	1 !
012	0010010000100000	il i	00100001000000001	li	000001000000000001	1	00000001000100001	1 ;
013 014	0001001000100000	il i	00010000100000000		01001000010000001	1	0100001000000000001	1 ;
014	0001000100010000	il i	00001000010000001	1	01000100001000001	1	00100001000000001	1 1
016	0000100100100000	i	00000100001000001	1	01000010000100001	1	00010000100000001	Ιí
017	10000100010010000	1 1	00000010000100001		00100100010000001		00000100001000001	
018	0000000100000000	11 1	001001001001000001	1	00100010001000001	Ιi	00000010000100001	. 1
019	0000000000100000	1 1	00100000100000001		00010010010000000	li	01001001001000001	1
020	0000000001000000	1 1	0001000001000000		00010001001000001		101001000100100001	1
021	0010010000010000	1 1	0000100000100001		[00010000100100001	1	010000010000000001	
022 023	0000100100010000	il î	00100000000100001		0000001000000000000	1	00100100100100001	1
024	0000010010010000		0010000000100000	1 1	01001000001000001	1 1	00100000100000001	li
025	0000000010000000	1 1	0001000000010000	1 1	01000100000100001		000010000010000001	li
026	100001000000000000	0 2	001001000000000010		0010010000100001		00000100000100001	Ī
027	0010010010000001	0 2	00010010000000010	2 2	0001001000100000	il i	01001001000100000	1
028	0010001001000001	0 2	00001001000000010		00010001000100001	il î	101000100100100001	1
029	0010000100100001	0 2	00000010010000010		[0000100100100000]	1 1	01000000100000000	
030 031	0010000010010001		00000001001000010	j 2	10000100010010000	1 1	00100000010000001	
032	10000010000000000000001	ol 2	00000000100100010) 2	000000100000000	11 }	00001000000100001	ili
033	10010010001000001	0 2	00000000010010010	2	0000000001000000		01000000000100001	il i
034	10010001000100001	0 2	0010001000000001	$\begin{bmatrix} 2 \\ 0 \end{bmatrix}$	0010010000010000	ilî	0100000000100000	1 1
035	0010000100010001	0 2	0001000100000001		0001001000010000	il ī	1001000000000100001	1 1
036 √	0010000010001001	0 2	0000010001000001		10000100100010000	וו	0100000001000000	1 1
037 038	0001001001000001	0 2	0000001000100001	0 2	10000010010010000	1 1	001000000100000	
039	0001000010010001		0000000100010001	0 2	0000000010000000		010010000000000001	il i
040	10001000001001001	0 2	00000000010001001	0 2	0100100100000001		00100100000000010	2 2
041	10000001000000001	0 2	0010000100000001	0 2	010001001000001	0 2	000100100000000010	0 2
042	0000000000100001		0001000010000001	0 2	0100000100100001		1000010010000000010	0 2
043	0010010000001001	0 2	0000010000100000		0100000010010001	012	1 0000010010000001	0 2
044	0001001000001001		0000001000010001		0100000001001001	0 2	0000001001000001	0 2
045	0000010010001001	ŏl ž	0000000100001001	0 2	inna1000000000001	0 2	0000000100100001	
046 047	0000001001001001		10010010010010001	0 2	0100100010000001	0 2	0000000010010001	
048	0000000001000001	0 2	0010010001001001	0 2	0100010001000001		01000100000000001	
049	10010010000010001	012	0010000010000001	0 2	010000100010001		00100010000000001	0 2
_050	0010001000001001	0 2	0001001001001001	0 2	0100000010001001	ol ž	10001000100000000	0 2
-051	0001001000010001	0 2	0001000001000001		0010010010000001	ol 2	10000100010000001	0 2
052	0001000100001001	0 2	100001000010001	0 2	- loo (0000001	012	0000010001000000	0 J Z
053	0000100100010001		0000001300001001	ol 2	0010000100100001	0 2	0000001000100001	0 2
054 055	0000010010010001	0 2	0010000000001001	0 2	Inn10000010010001	OI 2	0000000100010001	0 2
055 056	0000010001001001	ŏl ž	10010000000010001	012	Inn 10000001001001	0 2		
057	[00000000010000000]	0 2	100010000000001001	.012	00000100000000001 0100100001000001	0 2 0 2		O 2
058	10010010000100001	0 2	0010000000100001	0 2	0100010000100001	ol ž	10001000010000001	0 2
059	10010001000010001	0 2	0001000000010001		I AAAA 4 AAAA 1 AAA 1	Ŏ Ž		0 2
060	0010000100001001	0 2	Imminimization	<u> </u>	1010001000			

FIG. 12B

DATA	DCG11		1	DCG12		DCG21		DCG22		
SYMBOL	Code Word	NCC	Code	₩ord LSB	NCG	Code Word	NCG	Code Word	CD.	NCG
061	0001001000100001			100010010	2	MSB LSB 01000001000010010	2	00000100001000	SB 1010	2
062 063	0001000100010001			010010010	2	00100100010000010		00000010000100	010	2
064	0001000010001001		00100000	0010000010	2 2	00100010001000010 00100001000100010		000000010000010		2
065	0000100010010001		000010000		2	00100000100010010	2	01001000100100		2
066	0000100001001001		000001000		2	00010010010000010	2	01001000010010	010	2
067 068	1001000000000000100		001001000		3	00010001001000010	2	01000001000000		2
069	000100000000000000000000000000000000000		000100100		3	000100000100100010	2 2	/ 00100100100100100 00100100010010		2 2
070	000100000000000100) 3	000001001		ă	000000100000000010	2	00100000100000		2
071	0000100000000010		000000100	10000100	3	010010000000010010	2	00010010010010	010	2
072 073	00100100100000100	_	0000000010		3	001001000000010010	2	00010000010000		2
074	00100001001000100		001001000		3	00010010000010010 00001001000010010	2 2	00001000001000		2
075	00100000100100100) 3	001000100		3	00000100100010010	2	00000010000010		2
076	00010000000010000		000100100		3	00000010010010010	2	010000000000010		
077 078	1000010000000001000		000100010 000010010		3	00000000001000010 01001000000100010	2 2	010000000000100		2
079	00000000010010000		000010001		3	01000100000010010	2	0010000000000010		2 2 2 2
080	00000000001001000		000001001	00001000	3	00100100000100010	2	00100000000100		2
081 082	000000000100010000 00000000010001000		000001000 000000100		3	00100010000010010	2	00010000000010		2 2 2
083	00000000001000100		000000100		3	00010010000100010 00010001000010010	2 2	010010010000100 01000100100100		2
084	00000001000010000	3	000000010		3	00001001000100010	ž	01000010010010		2
085 086	000000000100001000		000000010		3	00001000100010010	2	010000000100000		2 2 2
087	00000000010000100 00100100010010000		001001000 001000100			00000100100100010 00000100010010010	2 2	00100000001000		2
088	00100100001001000		001000010			900000001000000101	2	000010000000100		2
089	00100100000100100	3	000100100	00010000	3	01001000001000010	2	010010010001000		2
090 091	00010010010010000 00010010001001000	3	000100010			010001000000100010	2	010010001000100		2
092	00010010000100100	3	0000100010			01000010000010010 00100100001000010	2	010001001001001000 010001000100100		2
093	00001001001001000	3	0000100010	00001000	3	00100010000100010	2	010000001000000	101	2
094 095	00001001000100100 00000100100100100	3	000010000			00100001000010010	2	001001001000100		2
095	00000010000010000	3	0000100010			0001001000100010 00010001000100010	2 2	001000100100100 001000000100000		2
097	00000001000001000	3	0000100000		3	00010000100010010	2	000100000010000		2
098	1000000000100000100	3]	0000010001			00001001001000010	2	000010000001000		2
099 11 100	00100100100010000 00100100010001000	3	0000010000 0000010000			00001000100100010 0100000000000010000		0000010000000100 0100100000000001		2
101	00100100001000100	3	0010000001			00100000000001000		0010010000000001		3
102	00100010010010000	3	0010000000			000100000000000100		0001001000000001		3
103 104	00100010001001001	3	0010000000 0010000010			01001001000000100/ 01000100100000100/		0000100100000001 0000010010000001		3
105	00010010010001000	_ =	0010000001			100010010000100		00000100100001		3
106	00010010001000100	3	001.0000000			1000001001000100		000000010010001		3
107	00010001001001000		0001000001			1000000100100100		000000001001001		3
108 109	00010001000100100		0001000000 0001000000			00100000000010000		0100100000000010 01000100000000001		3
	00001000100100100		0010000100			000010000000000100		0010010000000010		3
	00000100000010000		0010000010			1001001000001000		001000100000001		3
	00000010000001000		0010000001			01001000100000100	3 1	000100100000010		3
	00000001000000100 001001001000001000		0001000010 0001000001			1000100100001000		000100010000001 000010010000010		333333
115	00100100010000100	3	0001000000	1000100	3 (1000010010001000	3 (000010001000001	00	ž
	00100010010001000		0000100001		3 0	1000010001000100	3 (000001001000010	00	3
	001000100010001000 0010000100010001000		0000100000			01000001001001001000 010000010001001001		000001000100001 00000010010010		3
	00100001000100100		010010010			0100100100000100		00000010010010010		3
	00010010010000100		0010001000			0100010010000100		000000010010010		3

FIG. 12C

DATA	DCC 11		DCG12		DCG21		DCG22	
SYMBOL	Code Word	NCG	Code Word	NCG	Code Word	NCG	Code Word	NCG
121	00010001001000100	3	00100001000001000	3	00100001001000100	3	00000001000100100	3
122	00010000100100100	3	00100000100000100	3	00100000100100100	š	01001000000010000	3
123	00001000000010000	3	00010001000010000	13	00010000000010000	3	01000100000001000	3
124	00000100000001000	3	00010000100001000	3	00001000000001000	3	01000010000000100	3
125	00000010000000100	3	00010000010000100		000001000000000100	3	00100100000010000	
126	00100000000100000	4	00000000001000000	4	00000000010010000	3	00100010000001000	3
127 128	00000100100000000	4	00000000010000000	4	00000000100010000	3	00100001000000100	3
120	. 00000001001000000	4	001001001000000000	4	000000000010001000	3	00010010000010000	3
129 _t	00000000100100000	4	00100001001000000	4	00000000100001000	3	00010000100000100	3
131	00001000100000000	4	00100000100100000	4	00000000010000100	3	00001001000010000	3
132	00000100010000000	4	00100100010000000	4	01001000010010000	3	00001000100001000	3
133	00000010001000000	4	00100010001000000	4	01001000001001000	3	00001000010000100	3
134	00000001000100000	4	00100001000100000	4	01001000000100100	3	00000100100010000	3
135 136	00010000100000000	4	00010010010000000	4	00100100010010000	3	00000100010001000	3
137	00000100001000000	4	000100001001000000	4	00100100001001000	3	00000100001000100 0000001001001001000	3
138	00000010000100000	4	00100100001000000	4	00010010010010000	3	00000100001001000	3
139	00100100100100000	4	00100010000100000	4	00010010001001000	3	00000100000100100	3
140	00100000100000000	4	00010010001000000	4	00010010000100100	3	01000000010010000	3
141 142	0001000001000000	4	00010001000100000	4	00001001001001000	3	0100000000010010000	3
143	00000100000100000	4	00001001001000000	4	00001001000100100	3	010000000000100100 01000000100010000	3
144	00100000010000000	4	00100100000100000	4	00000010000010000	3	01000000010001000	3
145	00010000001000000	4	000100100000100000	4	0000000100001000	3	010000000001000100	3
146	00001000000100000	4	00001001000100000	4	00000000100000100	3	00100000010010000	3
147	00100000001000000	4	00000100100100000	4	01001000100010000	3	001000000001001000	3
148 149	00010000000100000 01000000010010001	4	00000000100000000 01001000010010001	1 1	01001000010001000 01001000001000100	3	001000000000100100	ភភភភភភ
150	01000000001001001		010010000001001001	i	010001000010010000	3	01000000100001000	ž
151	01000000100010001		01001000100010001	ī	01000100001001000	3	01000000010000100	3
152	01000000010001001		01001000010001001	1	01000100000100100	3	00100000100010000	3
153 154	01000001000010001		01000100010010001	1	00100100100010000	3	00100000010001000	3
155	01000000100001001 01001001001001001		01000100001001001 010000000000010001	1	00100100010001000 0010010001000100	3	00100000001001001	3
156 5	01000010000010001		01001001000001001	î	00100010010010000	ã l	00010000001001000	3
157	01000001000001001		01000100100001001	1	00100010001001000	3	00010000000100100	3
158	010010000000001001		01000010010001001	1	00100010000100100	3	01001001001001000	3
159 -⊼ 160	01001000000010001 010001000000001001		01000001001001001001 01001001000010001	1	00010010010001000	3	01001001000100100 01000100100100100	3
161	010001000000010001		01001000100001001	1	00010010001000100		01000010000010000	š
162	01000010000001001		01000100100010001	i l	00010001000100100		01000001000001000	3
163	01001001000000001	1 [0	010001000000000001	1	00001001001000100	3	01000000100000100	3
164	01000100100000001		010000100000000001	1	00001000100100100		00100001000010000	3
165	01000010010000001		01001001001000001	1	00000100000010000		001000001000010001	3
166 167	01000001001000001		01000001000000001	1	00000010000001000		001000000100001000	3
168	01001000100000001		1001001000100001	i I	01001001000010000		00010000010001000	3
169	01000100010000001		1000100100100001	i	01001000100001000	3 1	00010000001000100	3
170	01000010001000001		1000000100000001	1	01001000010000100		00001000010010000	3
	01000001000100001	1 19	010000000000100001	1	01000100100010000		00001000001001000	3
172 173	01001000010000001		01000000001000001	1	01000100010001000		00001000000100100	3
174	01001001000000010	2 1	100100000000000010		0100001001001000100		01001000100100100	3
175	01000100100000010		1000100000000010		01000010001001000		01000100000010000	3
176	01000010010000010	2 0	1000010000000010	2	01000010000100100	3	01000010000001000	3
	01000001001000010		1001001001000010		00100100100001000		01000001000000100	3
	01000000100100010 01000000010010010		1001000100100010		0010010001000100	3	00100100100100100	3 3
		2 0	1000001000000010	2	00100010010001000 00100010001000100	3	00100001000001000	3
100	2237003000010	٧ -	101000000000000000000000000000000000000			~ 1	~100010001001	<u> </u>

FIG. 12D

SYMBOL NSS	DATA	DCG11		DCG12		DCG21		DCG22	
Table			NEC		NCC	Code Word	NCC	Code Word	NCC
183	101		1	1120 1'28					
184					2] 3		
1845					5				
185	184				2				
188					2	00010000100100100	3		
188					2				3
199					2				
191 01000100001000010 2 010000000000					2		_		
193 1000010000010000 3 100110000000100 3 100110000000100 3 1001100000000100 3 100110000000100 3 100110000000100 3 100010010000000 4 1010001001000000 4 1010001001000000 4 1010001001000000 4 1010001001000000 4 1010001001000000 4 1010000000000				01000100100100010	2				-
194 1010010100000001000 3 0100110000000100 3 000000001010100000 4 011000100000000 4 195 010001010100000100 3 01001100000001000 3 000000001001000000 4 00100010000000 4 197 01000001001001001 3 01001100000001000 3 0000000001001000000 4 0010001001000000 4 197 01000001001000000 3 0100010000000100 3 0000000000					2				
195 01000100100000100 3 010010000000100 3 0100100000000100 3 0100100000000100 3 0100100000000100 3 0100010000000100 3 010000000000									-
195 01000100100000100 3 01000100000000 3 0000001000100000 4 001000101000000 4 001000101000000 4 001000101000000 4 0010000101000000 4 0010000101000000 4 001000101000000 4 0010000101000000 4 0010000101000000 4 0010000101000000 4 0010000101000000 4 0010000100						-	- 1		
198 01000001001001001 3 0100010000001000 3 00000100001			3	01000100000000100	3				
198 01000001001001001 3 01000000100100 3 01000000100000 4 010010000000 4 01001001000000 4 010010000100			3						4
199					3				-
									-
202 0100100001001000					3		-		-
203 0100100000100100 3 010000001000100 3 010010010100000 4 001000100001000					3 1		4		
205									
205					3				
200	205								-
208				01000000100001000	3				4
200 0100010000010000 3 01000100100100 3 00110000011000000 4 0011001000100000 4 011001000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011000100000 4 0110011001000000 4 011001100100000 3 01100000011000100 3 0110000000000									-
210					3				-
211 0100100010000100 3 0100001000001 3 010000010000000 4 000010010000000 4 010001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 01001001000000 4 0100100100100000 4 010001001000000 4 010001001000000 4 010000001000000 4 010000000100000 4 010000000100000 4 010000000000									-
213			3	01000010000010000	3 [4
214									4
215									4
216									1
218			3	01000100000010000	3	010000000001001001	1	001000000000001001	ì
219									1
220									1
222	220								i
223									1
224									-
225									
227									-
228			4	01000001000100000			1 0		
229									
230 001000001001001 1 00000000100010			11	000000000000000000000000000000000000000					
231 000100001001001 1 0000001000010									
232 000100001001001 1 00000001000010	231		-				i là	0100100100001001	
234			- 1				1 0	00100010010001001)
235 00010000100010001 1 00010010010001 1									1
236 00010000010001001 1 0001001001001 1									
237 0000100001001001 1 00001001001001 1	236		i	00010010001001001					-
239 00100100000001001 1 0000000100001001 1 00001001			1 1	00001001001001001	1 0	0100100000001001	1 0		-
									-
									1

FIG. 12E

DATA	DCG11			DCG12		DCG	21		D	CG22	
SYMBOL	Code Word	NCG	Code NSB	Word LSB	NCG	Code Wo MSB	LSB	NCC	Code MSB	Word 1.SB	NCG
241	0000100100000100		00100100		1	00000010010		1	010010000		1
242	0000010010000100	1 1	00100010		1	00000001001		1 1	010010000		1
243	0001000010000100	1 1	00100010		ì	01001000000		1	001001000)
244	0000100010001000	1 1	00010010	010001001	1	01000100000		1	001001000)
245	0000100001000100	L 1	00010001	100100100	1	00100100000		1 1	000100100		1
246	0000010001001000	L) L	000001000	000010001	1	00100010000	001001	1 1	000100100		1
247	0000010000100100	1 1	000000100	000001001	1	00010010000	010001	1	000010010		1
248	0001000100000300	l I	001000000	000001001	1	00010001000	001001	1	000000100		1
249	0000100100001000	1 1	001000000	000010001	1	00001001000	010001	1	000000010		1
250	0000100010000100		000100000	000001001	1	00001000100	001001	1	010010001	00010001	11
251	0000010010001000		000100000	000010001	1	00000100100	010001	1	010010000	10001001	1
252	0000010001000100	l i	000010000	000001001	1	00000100010	001001	1	010001000	10010001	1
253	0000001001001000	il i	00100100	100001001	i	00000010010	010001	1	010001000	01001003	1
254	0000001000100100		001000100	10001001	1	00000010001	001001	1	001001001		1
255	0010001000001000	i i	001000010	1001001		01000100000	010001	1	001001000	10001001	1

FIG. 12F

DATA	DCG31		DCC32			DCG41		DCG42	
SYMBOL	Code Word	NCG	Code Word	NC	c	Code Word	NCC	Code Word	NCO
000	100100100000000001	-		2	4	MSB LSB	ļ	T NIZR TZR	
001	10001001000000001	l i	00100100000000000	- 1	ı	10010010000000001	1	00100100000000000	
002	10000100100000001	l i l	00001001000000000		1	1000100100000001	1	01000100000000000	
003	10000010010000001	l i i	00000100100000000		ı	10000100100000001]]	0010001000000000	
004	10000001001000001	l î l	0000001001000000	il i	1	10000001001000001	1	1000010000000000	
005	10000000100100001	i	00100010000000000	$i \mid i$		10000000100100001	1	0100001000000000	
006	10010001000000001	1	0001000100000000	ו ונו		10010001000000001	l i	10010010010000001	
007	10001000100000001)	00001000100000000	1 1		100010001000000001	i	10010001001000001	li
800	10000100010000001	1	0000010001000000			10000100010000001	1	10010000100100001	li
009 010	1000001000100001	1	0000001000100000			100000100010000001	1	100000100000000001	
011	1001000010001000001	1	10000100000000000			10000001000100001	1	01001001001000001	1
012	10001000010000001	il	0010000100000000			010010010000000001	1	01001000100100001	1
013	10000100001000001	î l	0000100001000000			01000100100000001 010000100100000001	1	01000001000000001	
014	10000010000100001	i	0000010000100000			1000001001000001	1	00100100100100001	1 1
015	00100100100000001	1	0000001000010000	īl ī		1000000100100001	î	1.0010010001000001	1 1
016	00100010010000001	1	1001001001000000	1 1		0010000100000001	ì	10010001000100001	li
017	00100001001000001	1	1001000100100000	1 1	þ	0001000010000001	ì	10001001001000001	Ιî
018 019	00100000100100001		1001000010010000		11	10000010000100001	1	10001000100100001	lī
020	10010000010000001	1 1	100000100000000000	1 1		0000010000100001	l	10000001000000001	1
021	10001000001000001	i f	0010010010010000 001000001000000			1001000100000001	ļ	0100100100010001	1
022	10000100000100001	\mathbf{i}	00010000010000000			10001000100000011 10000100010001000011	1	01000100100100001	!
023	00100100010000001		00001000001000001			1000001000100001	1	01000000100000001	1 1
024	00100010001000001	1	00000100000100001	il ī		0100100100000001	i	10000000000100001	1 ;
025	00100001000100001		1001001000100(1001			0100010010000001	i	100000000001000001	l i
026 027	00010010010000001		10010001000100001			0100001001000001	1	01000000000100001	ì
028	00010001001000001		10001001001000001 10001000100100100001			0100000100100001	1	10000000010000001	1
029	00000010000000001		10001000100100001			0010000010000001	1	01000000001000001	1
030	10010000001000001		0100000010000001			0001000001000001 0000100000100001	1 1	00100000000100001	1
031	10001000000100001		0010000001000001			1001000010000001	i	10001001000100001	i
032	00100100001000001	1 0	00001000000100001			1000100001000001	î	10000100100100001	lî
033	00100010000100001	j j	0000000000100001	1	30	0010000000100001	1	10000000100000001	Ĭ
034 035	00010010001000001	1 1	0000000001000001	1		1001000000100001	1	010000000010000001	1
036			.00000000010000001 0100000000100001			1100100000100001	1	001000000001000001	2
037			0010010000100001	1		0010010000000010 00010010000000010	2	010010000000000010 0010010000000000010	2
038		î lî	0001001000100001	l î		0000100100000010	2 2	100010000000000010	2 2
	00100100000100001	1 1	0000100100100001	lil		000010010000010	2	010001000000000010	2
			0000000100000001	1		000001001000010		001000100000000010	2
		jΩ	01000000001000001	1		000000100100010	2	10000100000000010	2
		2 0	0010000000100001 010010000000000010	1		000000010010010		010000100000000010	2
			001001000000000010	2 2		010001000000010 0010001000000010		001000010000000010	2
A : A			0001001000000000101					100100100100000010 100100010010000010	2 2
			00001001000000010	2		000010001000010		10010000100100010	2
047	10000001001000010		00000100100000010				_ 1	10010000010010010	2
048	10000000100100010	2 10	00010000000000010	2			_ [10000010000000010	2
	100000000010010010		1000100000000010	2		001001000000010	2	01001001001000010	2
	100100010000000010 2 100010001000000010 2		010001000000010			00010010000001010		1001000100100010	2
			0001000100000010	2				01001000010010010	2
	10000100010000010 2 10000010001000010 2		00001000100000010 00000100010000010					01000001000000010	2
	10000001000100010		00001000000000010					XX 100 100 100 100 100 10 XX 100 100 100 100 100 100 100 100 100 1	2
055	0000000100010010		100001000000010					0100000100000010	2
	00001000000000010 2	00	010000100000010					00000000000100010	2
	0010000100000010 2		001000010000010	2				1000000000010010	2
	10001000010000010 2		000100001000010				2 1	0000000001000010	2
	0000100001000010 2 0000010000100010 2		000010000100010			200001000010010		10000000000100010	2
, ,,,,	0000010000100010 2	10	0100100100000101	2 (11(00100010000000101 2	: 10	0100000000010010	2

FIG. 12G

DATA		DCG31			DCG32			DCG41		D	CG42	
SYMBOL	ISB Code	Word LSB	NCG	LISB Cod	e Word LS8	NCG	Code MSB	Word LSB	NCG	Code 1	Word LSB	NCG
061	10000001	000010010	2	1001000	1001000010	2	01000100	010000010	2	100100100	00010010	2
062	00100100	100000010		1001000	0100100010	2	01000010	0001000010		100010010	00010010	2
063		010000010		1001000	0010010010	2	01000001	1000100010	2	100001001	00010010	2
064		001000010		1000001	0000000010	2	01000000	100010010	2	100000100	10010010	2 2 2
065		100100010	2		0100100010	2	00100100	100000010	2	100000000	10000010	2
066		010010010	2		0010010010	2	00100010	XX 10000010		0100000000	01000010	2
067		000000010	2	0010000	01000000010	2	00100001	1001000010		0010000000	0100010	2
068		010010010	2	0001001	0010010010	2	00100000	100100010		1001001000	0100010	2
069		010000010	2	0001000	0010000010	2	00100000	010010010	2	1001000100	00010010	2
070	10001000	001000010	2		0001000010	2	10010010	010010010	2	1000100100	0100010	2
071	10000100	000100010	2	0000010	0000100010	2	10010000	010000010		1000100010	00010010	2
072	10000010	000010010	2	0000001	0000010010	2	10001000	001000010	2	1000010010	0100010	2 2 2 2
073		010000010	2	1000000	0000100010	2	10000100	0000100010	2	1000010001		2
074		001000010	2		0001000010		10000010	000010010		1000000010	00000010	2
075		000100010	2	0010000	0000010010		01001000	010000010		0100100100	00010010	2
076		100010010	2		0000010010		01000100	001000010		0100010010		222222
077		010000010	2		1000010010		10010000	000010010		0100001001		2
078		100010010	2		0100010010	2		000010010		0100000001		2
079		001000010	2		0010010010	2	00100100	000010010		0010000000		2
		100100010	2		0010000010	2		000100010		1001001000		2
081		10010010	2		0000100010 j	2		000010010		1001000100		2
082		000010010	2		0000010010	2		000100010		1001000010		2
		000010010	2		0000100010	2		000010010	2	1000100100		2
		000010010			1000010010			000100010		1000100010		2 2 2 2
		000010010			10000100010			000010010	2	1000100001		2
		100010010			0100010010	2		001000010	2	1000000100		2 2 2
	000000100				100100010	2		000100010	2	0100100100		2
	100100000				010010010			000010010	2	0100100010		Z
	100010000				100000010	2		001000010	2	0100010010		2
	001001000				001000010			000100010	2	0100010001		2
	001000100				0000100010	3		000010010	2	01000000010		2
	000100100				000010010			001000010	2	0010010010		2
	000100010				001000010	2		000100010	2	0010001001		2
	000010010				000100010	2		000010010	2	10010000001		2
	000010001				100010010			000000100	3	0100100000		3
	000001001		ź I		001000010	2		000000100	3	0010010000		3
	000001000 100100000				0100100010			100000100 010000100	3	1001000000		223333
	100100000				000000010			001000100	3	1000100000		3
	100001000				100010010			100100100	3	0100100000		3
								000010000	3	0100010000		3
	001001 00 0 00100 010 0				010010010			00001000	3	0010010000		3
	0010000100				0010000010			000001000	3	0010001000		3
	0001000100			00010000		2	TOTOTO	7000001000		10010001000		3

FIG. 12J

DATA	DCG31		DCG32		DCG41			CG42	
SYMBOL	Code Word NSB LSB	NCG	Code Word	NCG	Code Word MSB LSB	NCG	Code MSB	Word LSB	NCG
241 242 243 244 245 246 247 248 249	10000001000001001 00100000100010001 001000000	1 1 1 1 1 1 1	00100100100010001 0010010001001001001 0010001001	1 1 1 1 1 1 1 1 1 1	01001000000010001 0100100000001000 001001	1	MSB 01000100 01000010 10010001 10010001 10001000 10001000 100001000 100001000 010010	00001001 010001001 00001001 00001001 00010001 100010001 100100	1 1 1 1 1 1 1 1 1
251 252 253 254 255	00100000100001001 0001000010010001001 000100001001)]	100001001000001001 1000010010010001001 1000001001		01000100100001000 01000100010000100 0100001001	3	010010010 010010001 01001000 010001000	00001000 10000100 00010000	37777

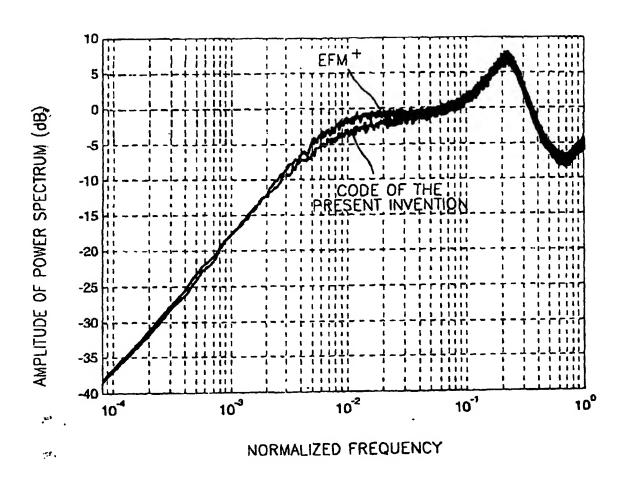
FIG. 13A

	ACGI		ACG2		VCC3		ACG4	
DATA SYMBOL	Code Word	NCG	Code Word	NCG	Code Word NSB LSB	NCG	Code Word MSB LSB	NCG
000 001 002 003 004 005 007 008 009 010 011 012 013 014 015 016 017 018 019 020 021 022 023 024 025 027 028 029 030 031 031 031 031 031 031 031 031 031	00000100000000 00001000000000 001001000000	4 4 4 4 4 1 3 1 2	000001000000000 000010000000000 00100100	4 4 4 4 2 2	100100010000001 10010010000001 100100100	111112222222233333333333333344444444441111111	000001000000001 00001000000001 000010010	111111111111111000000000000000000000000

FIG. 13B

DATA	<u></u>		VCC1				AÇC	2			ACG3			ACG4	
YMBOL	IISB	Code	Word	LSB	NCG	NSB Co.	de Vo	rd LSB	NCG	Code	Word LSB	NCG	Code NSB	Kord	NO
061 062													00010000	0001000	1 3
063	ļ				1 1				1			1 1	00010000	0010000] 3
064	1											1 1	00010010	0100100	1 3
065	l											1 1	00010010	0100000	1 4
066	l							- 1	· •			}	00000100	0000000	14
067	Ì				1			- 1	- 1			1 1	00001000	0000000	1
068					1			- 1	- 1			1 1	00000100	0100000	1 4
069		•	•	1	,			- 1	- 1			[00000100	1000000	1 4
070				- 1	- 1			- 1	- 1			1 1	00001000	0100000	1 4
071				- 1	- 1			- 1	- 1				00001000	1000000	4
072				- 1	- 1			1	- 1			. 1	00001001	9000000	4
073				- 1	,			1	- 1				000100000	0100000	4
074				ı	ł			ł	- 1			1	00010000	1000000	1 4
075				- 1	- 1			- 1	- 1			- }	000100011	0000000	4
076				ı	- 1			- 1				1	000100100		1 4
077				1	ı			ì	1		l	- 1	001001000		4
078					- 1			- !			- 1	1	01001000		4
079				- 1	- 1			ı	1	•	- 1	1	100001000		4
080				- I	- 1			- 1	- 1			- 1	100010000		1 1
081				- 1				1	- (- 1	100100100		4
082				- 1	- 1			- 1	- 1		- 1	- 1	000100000		li
083					1			- 1	ı		ĺ	I	000001001		li
084				- 1	- 1			- 1	- [- 1	ı	000010001	nninni	1
085					- 1			- 1	- 1		•	- 1	000010010	001001	i
086					- 1			}	- 1		1	ſ	000010010		i
087 088				- 1	- 1			j	- 1		1	- 1	000100001	001001	ī
089				- 1	- 1)	1		1		000100010	001001	1
090				- 1	- 1			- 1	- 1		- 1	1	000100010	010001	1
91				ł	- 1				- 1		· 1		000100100		1
92				- 1	ı			1	1		}		000100100		1
žěš l				- 1	- 1			- [- 1		1		000001000		1
94				1	- }			1	- 1		!	- 1	000001000	010001	1
95 3				- 1	1			- 1	- 1		1		000010000		1
96				- 1	- 1			- 1	- 1		j		000010000		1
97								- [1		1		0001000000		1
÷55.													1001001000	1.00071	ı

FIG. 14





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(12)

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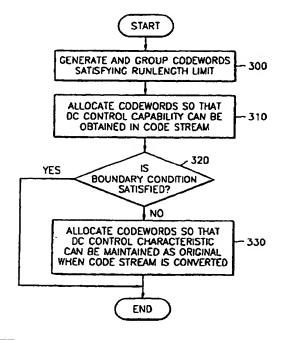
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(54) Code generation and allocation method

(57)A method for generating and allocating codewords is provided. The method includes allocating one of two selectable codewords b1 and b2 as codeword b when a preceding codeword a and a following codeword b form a code stream X, in which codewords b1 and b2 have opposite INVs which are parameters indicating whether the number of '1s' contained in a codeword is an odd number or an even number and when the code stream of a and b1 is X1, and the code stream of a and b2 is X2, allocating codewords such that the INVs of X1 and X2 are maintained to be opposite when a or b1(b2) should be replaced by another codewords in compliance with a predetermined boundary condition given between codewords. According to the method, by using a short codeword having less bits as a main conversion codeword, high efficiency is achieved in recording density. Also, when codewords which do not satisfy the run length conditions are replaced by other codewords, the codewords are allocated so that the DC suppression capability of the code stream can be maintained, and therefore higher DC suppression capability of the code stream is provided.

FIG. 3



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EUROPEAN SEARCH REPORT

Application Number EP 02 25 2790

	OCUMENTS CONSIDER Citation of document with indice			elevant	CLASSIFICATION OF THE
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ANNEX TO THE EUROPEAN SEARCH REPORT ON EUROPEAN PATENT APPLICATION NO.

EP 02 25 2790

This annex lists the patent family members relating to the patent documents cited in the above-mentioned European search report. The members are as contained in the European Patent Office EDP file on The European Patent Office is in no way liable for these particulars which are merely given for the purpose of information.

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